

ABSTRACT

Yoon et al. proposed a new efficient remote user authentication scheme using smart cards in 2004. This paper analyzes the proposed scheme and then proves that the password change phase of Yoon et al's scheme is still insecure.

KEYWORDS: Cryptography, Authentication, Smart cards, Password, Security attacks.

I. INTRODUCTION

Lamport [17] proposed the first well-known remote password authentication scheme using smart cards. In Lamport's scheme, the AS stores a password table at the server to check the validity of the login request made by the user. However, high hash overhead and the necessity for password resetting decrease the suitability and practical ability of Lamport's scheme. In addition, the Lamport scheme is vulnerable to a small n attack [7]. Since then, many similar schemes [23]-[26] have been proposed. They all have a common feature: *a verification password table should be securely stored in the AS*. Actually, this property is a disadvantage for the security point of view. So far different types of password authentication schemes with smart cards [3] - [4] - [5] - [6] - [12] - [13] - [14] - [18] - [20] - [21] - [24] - [29] have been proposed. In 2002, Chien-Jan-Tseng [13] introduced an efficient remote user authentication scheme using smart cards. In 2004, Ku and Chen [31] pointed out some attacks [7]-[28]-[30] on Chien - Jan and Tseng's scheme. According to Ku and Chen, Chien et al.'s scheme is vulnerable to a reflection attack [7] and an insider attack [30]. Ku and Chen claimed that Chien et al.'s scheme is also not repairable [28]. In addition, they also proposed an improved scheme to prevent the attacks: reflection attack and an insider attack on Chien-Jan-Tseng's scheme. In the same year, Hsu [10] pointed out that the Chien-Jan-Tseng's scheme is still vulnerable to a parallel session attack and Yoon et al. [11] claimed that the password change phase of improved scheme of Chien-Jan-Tseng's scheme is still insecure. This paper analyzes that the password change phase of Yoon et al.'s scheme is still insecure. The rest of the paper organized as follows.

Section II reviews Yoon et al.'s scheme [11]. Section III is about our observations on the security pitfalls in the password change phase of Yoon et al.'s scheme. Finally, comes to a conclusion in the section IV.

II. YOON ET AL.'S SCHEME

This section briefly describes Yoon et al.'s scheme [11]. This scheme also has four phases: the registration phase, login phase, verification phase and the password change phase. All these four phases are described below.

A. Registration Phase

This phase is invoked whenever U initially or re-registers to AS. Let n denotes the number of times U re-registers to AS. The following steps are involved in this phase.

- ❖ User U selects a random number b and computes $PW_S = f(b \oplus PW)$ and submits her/his identity ID and PW_S to the AS through a secure channel.
- ❖ AS computes two secret numbers $V = f(EID \oplus x)$ and $R = f(EID \oplus x) \oplus PW_S$, where $EID = (ID || n)$ and creates an entry for the user U in his account database and stores $n = 0$ for initial registration, otherwise set $n = n + 1$, and n denotes the present registration.
- ❖ AS provides a smart card to the user U through a secure channel. The smart card contains two secret numbers V, R and a one-way function f .

- ❖ User U enters her/his random number b into his smart card.

B. Login Phase

For login, the user U inserts her/his smart card to the smart card reader and then keys the identity and the password to gain access services. The smart card will perform the following operations:

- ❖ Computes $C_1 = R \oplus f(b \oplus PW)$ and $C_2 = f(C_1 \oplus T_U)$. Here T_U denotes the current date and time of the smart card reader.
- ❖ Sends a login request $C = (ID, C_2, T_U)$ to the AS.

C. Verification Phase

Assume AS receives the message C at time T_S , where T_S is the current date and time at AS. Then the AS takes the following actions:

- ❖ If the identity ID and the time T_U is invalid *i.e.* $T_U = T_S$, then AS will rejects this login request.
- ❖ Checks, if $C_2 = f(f(EID \oplus x) \oplus T_U)$, then the AS accepts the login request and computes $C_3 = f(f(EID \oplus x) \oplus T_S)$. Otherwise, the login request C will be rejected.
- ❖ AS sends the pair T_S and C_3 to the user U for mutual authentication.
- ❖ If the time T_S is invalid *i.e.* $T_U = T_S$, then U terminates the session. Otherwise, U verifies the equation $C_3 = f(C_1 \oplus T_S)$ to authenticates AS.

D. Password Change Phase

This phase is invoked whenever U wants to change his password PW with a new one, say PW_{new} . This phase has the following steps.

- ❖ U inserts her/his smart card to the smart card reader and then keys her/his identity and the old password PW and then requests to change the password.
- ❖ U 's smart cards computes $V^* = R \oplus f(b \oplus PW)$.
- ❖ Compare this calculated value V^* with the secret value V , which is stored in the smart card of the user U . If they are equal, then U can select a new password PW_{new} , otherwise the smart card rejects the password change request.
- ❖ U 's smart cards computes a new secret number $R_{new} = V^* \oplus f(b \oplus PW_{new})$ and then replaces R with R_{new} .

III. SECURITY ANALYSIS OF THE PASSWORD CHANGE PHASE OF YOON ET AL.'S SCHEME

Although, the password change phase of Ku and Chen's scheme is modified by Yoon et al. [11] to remove its security weaknesses. But, we analyze that the modified password phase of Yoon et al.' scheme is still not secure. This section discusses the security weaknesses of the password change phase of Yoon et al.'s scheme and proves that the modified phase is still vulnerable to security attack.

A. Security weaknesses in the Password Change Phase against the Outsiders

Observe the password change phase of Yoon et al.'s scheme, to replace/change the old password PW with a new password PW_{new} , the user/performer should be in possession of the old password PW . The following section describes how any outsider /malicious user can recover the password PW first and then apply this piece of information to make for the success of her/his attack.

It is clear that the smart card of a legal user U in Yoon et al.'s scheme contains: *the secret value* V , R , and *a random number* b and *a public hash function* f . According to Kocher et al. [22] and Messerges et al. [29], for the security point of view, to store the secret information in smart cards is not a good practice. On the basis of these assumptions [22]-[29], an antagonist is able to breach the secrets V , R and b , which are stored in the smart card of the user and then he will be able to perform a password guessing attack to obtain the password. For the success of this attack, by using the breached secrets R and b the adversary will perform the following operations:

- *The antagonist intercepts the login request* $C = (ID, C_2, T_U)$ and guesses a password PW^* .
- *Computes* $C_1^* = R \oplus f(b \oplus PW^*) = f(ID \oplus x)^*$ and $C_2^* = f(C_1^* \oplus T_U)$.
- *Checks if* $C_2^* = C_2$, then the adversary has correctly guessed the password $PW^* = PW$ and $C_1^* = C_1$. Otherwise, the adversary goes to step: 1.

Once the adversary has correctly obtained C_1 , instantly, the password PW^* corresponding to C_1 will be the correct password and then successfully, he can change the password of the user U . Consequently, when the smart card was stolen, the antagonist is able to recover the password PW of the user and once the adversary has correctly obtain the password PW , then he will be able to destruct anything of his choice. Since our focus and aim is to show that the password change phase of Yoon et al.'s scheme, which is shown below that an authorized user (antagonist) can easily replace the old password PW by a new password of her/his choice. For the success, the antagonist applies the following actions.

- Inters the smart card into the smart card reader, enters the identity ID and any password PW and then requests to change the password.
- The smart card of the user computes $V^* = R \oplus f(b \oplus PW)$ and then compare the computed value V^* with the stored value V . Obviously, both the value will be the same, because the adversary has entered the correct password. In this way, the smart card accepts the password change request.
- Selects a new password PW_{new}^* and supplies it to the smart card reader and ultimately the smart card computes a new $R_{new}^* = R \oplus f(b \oplus PW) \oplus f(b \oplus PW_{new}^*)$ and then replaces R with R_{new}^* .

Thus, if the malicious user stole the user U 's smart card she/he will be able to make a destructive action of her/his choice. Thus, the adversary is able to change the password with a new password of his/his choice. Now the registered/ legal user U also will not be able to make a valid login request with her/his valid smart card because now the her/his old password PW will not work .

B. Security weaknesses in the Password Change Phase against the Insider

This section proves that the password change phase of Yoon et al.'s scheme is not secure against an antagonist insider at AS. In Yoon et al.'s scheme, observe the registration phase, the User U selects a random number b and computes $PW_S = f(b \oplus PW)$ and submits her/his identity ID and PW_S to the AS through a secure channel. It means the insider of AS is in possession of the number $PW_S = f(b \oplus PW)$ for the legal user U . Again the AS computes two secret numbers $V = f(EID \oplus x)$ and $R = f(EID \oplus x) \oplus PW_S$, where $EID = (ID || n)$. Thus, the insider of AS is also in possession of the secret numbers V and R for the legal user U .

Suppose the user U is using the same password PW continuously, which is supplied by the AS at the time of registration, then the insider at AS will be able to change the password PW with a new password of her/his choice. If the smart card is in possession of an antagonist insider at AS for short time, then first, the insider inters the smart card into the smart card reader and can directly supply the value V to the smart card reader. Either, he directly supplies V or in place of $f(b \oplus PW)$, he supplies the value PW_S without using the hash button. Next, the antagonist insider enters a new password PW_{new}^* and then the smart card computes a new $R_{new} = R \oplus f(b \oplus PW) \oplus f(b \oplus PW_{new}^*)$ and then replaces R with R_{new} .

Thus, if the malicious insider stole the user U 's smart card once, only for a small time and then he can replace the user's password forever in such a way that the user U also will not be able to make a valid login request with her/his valid smart card because now the her/his old password PW will not work properly.

Thus the Yoon et al.'s password change phase is still insecure and that is under the threat of poor reparability.

IV. CONCLUSION

This paper analyzed that security weaknesses still exist in the password change phase of modified scheme of Yoon et al.'s scheme. The password change phase is still vulnerable to security attacks by an outsider as well as an antagonist insider at AS. Thus, the security pitfalls still exist in Yoon et al.'s scheme.

V. REFERENCES

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